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#### SOUTHEAST UNIVERSITY



Looking before Crossing: An Optimal Algorithm to Minimize UAV Energy by Speed Scheduling with a Practical Flight Energy Model

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#### Background

Problem Modeling

Solutions

Simulation







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# UAV data collection application scenarios

GN

• A UAV collects data from wireless sensors or IoT devices (GNs) deployed along a straight line

AV

power line

- 1. power transmission line
- 2. water/oil/gas pipe
- 3. river/coast
- 4. road
- 5. ...

river

# **Energy consumption problem of UAVs**



- Key issues: limited energy supply on board.
- Flight power 1000 times communication power

Flight~400WWe focus onCommunication~300mWflight energy

• Related flight energy consumption models



distance-related energy model duration-related energy model

Too simple to be accurate

# **Our practical flight energy model**

• We conduct a set of real-world flight tests



• The flight power is a convex function of the flight speed

#### **Consistent with theoretical analysis**

• Consistent with a most recent theoretical analysis.



UAV speed V (m/s)

### **Challenges to our problem**



1. min UAV flight energy to collection all data



2. GNs compete for UAV time to upload data

Three GNs compete to upload their own data



Such competition is complicated: each GN has

- 1. a different amount of data
- 2. a different transmission range size

Best GN transmission scheduling must be found





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#### **GNs and UAV**



- m Ground Nodes (GNs) are deployed alone a line
- The UAV flies straightly at a fixed height
- The UAV collects data from GN when flying



# Transmission range and required time



- GN *i* has a transmission range  $(s_i, f_i)$
- Within range, GN *i* requires  $\tau_i$  time to upload data
- We allow the ranges be different but aligned  $-0 = s_1 < s_2 < \cdots < s_n$  $-0 < f_1 < f_2 < \cdots < f_n = D$



#### **UAV speed scheduling**

- Speed scheduling function v(t)
  - The speed at time *t*
- Position of the UAV at time t

$$d(t) = \int_0^t v(\tau) \mathrm{d}\tau$$



#### **Range/completion constraint**



- Transmission ranges are overlapped but aligned, so the UAV collects data following GN index.
- At switching time  $t_i$ : the UAV finishes collecting data from GN *i* and starts GN *i* + 1
- Range constraint

$$s_i \le d(t_{i-1}) < d(t_i) \le f_i$$

Completion constraint

$$t_i - t_{i-1} \ge \tau_i$$

# **USS-GTS problem**

- Flight power p(v) for UAV flight speed v
- The energy consumption of UAV  $E = \int_{t}^{t_n} p(v(t)) dt$ Power(W) 400
- USS-GTS problem: find speed scheduling function v(t) to
  - 1. minimize UAV energy consumption E
  - 2. satisfy range constraint
  - satisfy completion constraint 3.



 $\frac{10}{\text{Speed}(\text{m/s})}$ 

15

550

350





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# **Distance accumulation trajectory**



- Time-distance diagram: reaching position d at time t
- distance accumulation function d(t) a curve
- speed scheduling function v(t) is the curve's slope





- Trajectory does not go freely on the diagram
- *completion constraint:* temporal on transmission times
- range constraint: spatial on transmission range





• For GN *i*, we draw a rectangle Room *i*: the north wall at  $d = f_i$ , the south wall at  $d = s_i$ , the west wall at  $t = t_{i-1}$ , the east wall at  $t = t_i$ 





- *Crossing-the-rooms* problem essentially asks:
  - 1. how to determine the length of each room
  - 2. how to design the trajectory crossing all the rooms





- *Crossing-the-rooms* problem essentially asks:
  - 1. how to determine the length of each room
  - 2. how to design the trajectory crossing all the rooms equivalent to ask the GN transmission switching times equivalent to ask the UAV speed scheduling function

Therefore, the solution is uniquely mapped to the solution for the original USS-GTS problem.

#### Theorem: straight trajectory saves energy



• **Theorem**. *The optimal trajectory is straight between any two points, as long as this is feasible.* 



#### The proposed algorithm



Algorithm 2: USS-GTS-GENERAL

```
1 k = 0, d = 0, t_0 = 0;
2 s_{n+1} = f_n // dummy for loop purpose
3 while k < n do
       v_n^v = \infty, v_s^v = 0;
       for j = k + 1 to n do
 5
            t_i = t_{i-1} + \tau_i, v_n^d = (f_i - d)/(t_i - t_k),
 6
              v_s^d = (s_{j+1} - d)/(t_j - t_k);
            if v_s^d > v_n^v then
7
                 v_m = v_n^v, \, k_m = k_n^v, \, d_m = f_{k_n^v};
8
                 break:
 9
            else if v_n^d < v_s^v then
10
                 v_m = v_s^v, \, k_m = k_s^v, \, d_m = s_{k^v+1};
11
                 break:
12
            end
13
            if v_n^v > v_n^d then v_n^v = v_n^d, k_n^v = j;
14
            if v_s^v < v_s^d then v_s^v = v_s^d, k_s^v = j;
15
16
        end
        if v_n^v == v_s^v then v_m = v_n^v, k_m = n, d_m = f_n:
17
        if v_m > v^* then
18
            x = d:
19
            for i = k + 1 to k_m do
20
                 t_i = \max\{(s_i - x)/v^*, \tau_i\} + t_{i-1};
21
                x = x + (t_i - t_{i-1})v^*
22
23
            end
        end
24
        Connect (t_k, f_k) and (t_{k_m}, d_m);
25
        d = d_m, k = k_m;
26
27 end
```

# Looking before crossing rooms algorithm



The algorithm works in three phases

1. Construct rooms

-Room *i* with length  $\tau_i$ , the minimal required time

- 2. Find a walking trajectory across rooms
- 3. Some room lengths are enlarged





- At origin, the view area through door 1 is in pink and view area through door 2 is in red.
- No view through door 3, which is beyond the northern boundary
- We walk along the northern boundary, reaching the farthest doorjamb.





- At the new position, two new view areas are in pink and red respectively.
- Since the northeast corner is beyond southern boundary of the current view area
- We walk along the southern boundary to the farthest doorjamb.





- In room 3 and 4, trajectory slope larger than  $v^*$
- **Theorem.** Any point of the optimal trajectory has a slope no larger than  $v^*$ .
- We enlarge the lengths of room 3 and 4 to reduce the slope to  $v^*$





- All room lengths and trajectory are determined
- Theorem. Algorithm USS-GTS-GENERAL produces the optimal distance accumulation trajectory for the offline USS-GTS problem within  $O(n^2)$  steps.









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# **Online heuristic algorithm**



- Previous offline algorithm requires all information
- In the online algorithm, we assume the information of a GN can be obtained only if the UAV flies close

Online heuristic algorithm outline:



#### **Simulation Results**

• Performance comparison between online heuristic and offline optimal. imes10<sup>5</sup> 12 Energy consumption (J) online heuristic

offline optimal

In all three settings, online heuristic performances near optimal.



4













Simulation



## Conclusion



- 1. Investigate a UAV data collection problem from GNs deployed along a straight line
- 2. Real-world flight tests: a speed-related energy model
- 3. Propose the *looking before crossing* algorithm on time-distance diagram: the optimal offline solution
- 4. Present an online heuristic algorithm which performances near optimal.
- 5. Our study on the practical flight energy model and speed scheduling have shed light on a new direction on UAV-aided wireless communication.



# Thank You!

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